

## Outline

- CS270 recap: what's still a black box?
- Moore's law for ever
- The "walls" and surmounting them
- Go for Speed
- Bandwidth/Memory Wall
- The decade of ILP and frequency
- Power \& Energy wall \& the rise of multi-core
- Utilization wall - next level or game over?

Back to the present - PA6 details and issues
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## What was brushed under the carpet in CS 270

Combinational circuits are instantaneous
Minimalist vs Efficient
■ LC-3 can execute any program
$■$ But does it do it efficiently?
How fast can the machine go?
How much power does it consume?
What is the manufacturing cost?
$\square$ Economies of scale

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## Moore's Law

- Originally "formulated" by Gordon Moore (~1965) but other earlier observations too
- The number of transistors that can be inexpensively placed on an integrated circuit is increasing exponentially, doubling approximately every years [later amended to 2 years]
- http://en.wikipedia.org/wiki/Moore's_law

It has held true till now and is expected to hold for about 10 more years

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## Moore's Law corollaries

- For chip designers:
- "we had better ensure that the exponential growth is maintained" - or else the competition will $)^{-}$
- Main reason why the law is being sustained
- Other features of semiconductors \& computing technology are also growing exponentially (but at different rates):
- Frequency
- Die size
- Memory: density, and speed (bandwidth and latency)
- Hard drive and I/O devices (both capacity and speed)
- Networks
- Power
- Manufacturing cost

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## Moore's law of expectations

Better, faster, cheaper, lighter, ...
$\square$ We as computer scientists are providing the technology that is changing the world exponentially

- Hard challengesExciting potential
Always room to innovate - if you stop learning you stagnate


## Exponential growth implies

When two quantities grow exponentially, but at different rates, their ratio also grows exponentially. Consider,

$$
\begin{aligned}
y_{1} & =a^{x}, \\
\text { and } y_{2} & =b^{x} \quad \text { for } a \geq b \geq 1
\end{aligned}
$$

$$
y=\frac{y_{1}}{y_{2}}=\left(\frac{a}{b}\right)^{x}=\alpha^{x} \text { for } \alpha \geq 1
$$

## Moore's law walls

- Memory gap/wall:
- Memory bandwidth grows much more slowly than processor speeds (since mid 80s, this was addressed by ever increasing on-chip caches).
- ILP (instruction level parallelism) wall
- One way to exploit increased clock frequency was to increase the instruction-level parallelism on chip (deeply pipelined, out-of-order, VLIW, etc.) leading to complicated control logic.
- Power wall
- Power dissipation ability is also increasing exponentially, but at such a slow rate that it has effectively peaked.
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- Utilization wall
- Multi- and many-core trend cannot be sustained: there is no way to keep all the transistors on future chips active at all times, and this "dark silicon" will also be increasing (exponentially)


## Improving the speed

Increasing the frequency

- Simply riding More's law

Better architecture through instruction level
parallelism (ILP)

- Pipelining
- Super-scalar

■ Out-of-order execution

## Tackling the memory wall

Caches (since mid eighties)
$\square$ If access to memory is slow, then build a faster (smaller) memory on-chip
Need to exploit
■ Locality of reference

- Reuse of data
- Collaboration between architecture, compiler and operating system


## Tackling the power wall

- Increasing the frequency implies increasing the heat generated
- And that has to be dissipated (or else th chip will melt
If you can't increase the frequency (raw speed)
- Add more processors (cores)
- "The processor is the new transistor" in Moore's law


## Back to the present (PA6)

Revisiting many of the concepts seen earlier:

■ HW4, LC3Viz and the "cycle-by-cycle" details of the instruction execution

- Appendix A of the textbook

■ PA2 (32-bit floating point addition) in C
■ PA5 (16-bit FPADD) in assembly

- Code comprehension

■ Reading code written by others

## Revisit normalization

- In scientific notation with 4-digit magnitude and l-digit exponent:
■ What is the largest number:

> - 9.999E9

- What is the smallest (positive) number?
= 1.000E-9

Can we do better?

- Don't force numbers to be (always) normalized:
- 0.001E-9

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## For binary numbers

Deal with the "implicit l"
■ Sometimes it's there and sometimes it isn't

- When?
- If exponent is 00000, no implicit l
- Need to handle corner cases in the algorithm.
- What if the inputs are a mix of normal and sub-normal numbers
- What if the answer is?

