# CS270 Programming Assignment 6 "LC-3 Lite Simulator (C Programming)"

Program due Saturday, December 10th, 2011 (via Checkin by 4:59pm)

No late submissions are allowed for this assignment.

Assignments are to be completed individually. Discussing implementation specifics of the program is **forbidden**.

Last updated: 12/06/2011 for due date extension/clarification (submit via Checkin, not RamCT)

# Goals

- Improve your C language skills (learn string and text file processing)
- Improve your understanding of the LC-3 machine architecture by building a simulator.

# The Assignment

In this assignment, you will develop a "lite" version of the LC-3 simulator using C programming. Due to the scale of this project, we provided you with a large portion of the program code and a working version of the simulator executable. The simulator is an extension of the basic LC-3 machine structure you implemented in PA2 (the files provided in this project include the solution for PA2). The executable will only work on the CS department lab machines (i.e., 64-bit HP Workstations). Also, the program has not been rigorously tested, so there might be a few minor issues.

Download the executable and skeleton code from here: <u>http://www.cs.colostate.edu/~cs270/Assignments/PA6/PA6.PROVIDED.tar.gz</u>

## **Running the Executable:**

The executable is named lc3sim-lite. It accepts a commands file similar to lc3sim. Many of the commands supported by lc3sim are supported by this version except breakpoint manipulation and single stepping through the program. The simulator itself supports most of the LC-3 opcodes except RTI and TRAP (will bail-out on assembly programs that have these instructions). A sample test case is provided for you to test it out.

\$ tar -zxf PA6.PROVIDED.tar.gz; cd PA6.PROVIDED

\$ cd testcases; ../lc3sim-lite cmds.shifttest.txt

# Part-A:

In the first part of the assignment, you will need to implement string tokeniznum\_tokensation and parsing the LC-3 symbol table file. The function implementations are located in infilesparser.c (and the declarations are in infilesparser.h). The function signatures are as follows:
/\*\*

\* Function to tokenize an input line into separate tokens

\* The first argument is the line to be tokenized and the second argument points to

\* a 2-dimensional array of char (i.e., an array of strings). The number of rows of this array should be

\* at least MAX\_TOKENS\_PER\_LINE size, and the number of columns (i.e., length

\* on spaces and tabs.

\*/

<sup>\*</sup> of each string should be at least MAX\_TOKEN\_SIZE). The tokens are delimited on

<sup>\*</sup> Returns 0 on success and negative value on failure

static int \_\_tokenize(char \*line, char tokens[][MAX\_TOKEN\_SIZE], int \*num\_tokens);
/\*\*
 \* Get the next symbol and its associated address from the symbol table file

- \* First arg is a file stream pointer to the symbol table file, and
- \* second arg is a 2-d array to hold the label string and the address string
- \*
- \* Returns 0 on success (and copies the symbol information to the 2nd arg),
- \* EOF if the file has reached end-of-file or a negative value on failure \*/

int symbol\_table\_getnext(FILE \*symfile, char label2addr[][MAX\_TOKEN\_SIZE]);

The tokenizing function splits the provided input string (argument-1) into separate tokens and stores those tokens in tokens[][] array. **The tokens are separated on space and tab characters**. The function must also trim any newline characters (or carriage return) and **must ignore empty lines**. The tokenizing function also sets num\_tokens to the number of tokens found.

The symbol\_table\_getnext() function reads the contents of a symbol table file and returns the next label and its associated address in the second argument label2addr[][]. That is, label2addr[0] holds the label string and label2addr[1] holds the address string. To understand how to parse the symbol table file, take a look at any LC-3 symbol table file from your previous assignments. Note that this function internally calls \_\_tokenize() to tokenize the input from the symbol table file. Look at the function commands\_file\_getnext() for an example on how to use the tokenizing function.

#### Note:

If you are unable to complete this part, you can proceed to the next part by linking your program against infilesparser.o.PROVIDED. In order to do that you need to modify the provided Makefile (the details are for you to determine). However you will lose points for not implementing Part-A.

## Part-B:

In this part of the assignment, you will implement the LC-3 machine in lc3sim.c. The main program calls lc3\_evaluate\_instruction(), which fetches the next instruction, decodes it and executes the instruction. The function returns with an error if it encounters the reserved instruction or an RTI instruction. **You do not need to implement the RTI and TRAP instructions**. The function returns a special LC3\_SIM\_HALT value upon executing the halt instruction (which is really the TRAP instruction with trap vector x25). As a start, we implemented the NOT, ADD, and AND instructions and provided a general structure for lc3\_evaluate\_instruction()

Note that the file lc3sim.h declares several functions. However, you are free to change these functions and implement the functionality in your own way (with certain rules described below). These are provided as a guideline for you to implement the simulated machine (our implementation is based on these functions).

## About Registers & Memory:

You do not have direct access to the LC-3 registers or the memory layout of the machine in the simulator. These are hidden away from you in the provided library file lc3memlib.a. However, we provided accessor functions to read/write from/to registers and memory. These are declared in lc3registers.h file. The description of these functions is as follows:

/\*\* \* Writes the contents of (argument) "word" to the register specified \* by "register\_name". Refer to Ic3machine.h for the register names. \* Prints an error if an incorrect register name is specified \*/ void write\_register(lc3\_register\_names\_t register\_name, lc3\_word\_t word); /\*\* \* Reads the contents of a register specified by argument "register\_name" \* and returns the value inside that register. \* Returns a null word if an incorrect register name is specified \*/ lc3\_register\_t read\_register(lc3\_register\_names\_t register\_name); /\*\* \* Writes the contents of MDR register into the location pointed to by MAR \* Aborts the program (and prints an error) if unable to write to a location \*/ void write memory(); /\*\* \* Reads contents of the location pointed to by MAR and stores it into MDR \*

\* Prints an error message to the console and sets MDR to null word if the \* location pointed to by MAR is not a location that has been previously

About Status Bits:

void read memory();

\*/

\* written into (i.e., an invalid read).

that the address to access is in MAR.

The LC-3 machine status bits can be accessed using the following accessor functions.

Refer to the file lc3machine.h for the register names. Note that read memory() and

write\_memory() function read the address from the MAR (memory address register) and read-from and write-to MDR (memory data register). Before calling one of these functions, you need to make sure

/\*\*
 \* Sets (to 1) the status bit specified by the argument "code" in the PSR.
 \* Refer to lc3machines.h for the status codes.
 \*/
void set\_status\_code(lc3\_status\_codes\_t code);
/\*\*
 \* Resets (to 0) the status bit specified by the argument "code" in the PSR
 \*/
void reset\_status\_code(lc3\_status\_codes\_t code);

/\*\*
 \* Tests whether a particular status bit (specified by the argument "code") is
 \* set to 1 or 0
 \*
 \* Returns 1 if the status bit is 1 in the PSR, otherwise returns 0
 \*/
int test status code(lc3 status codes t code);

Refer to the file lc3machine.h for the status code names. Note that our simulator does not implement other status codes defined by the LC-3 machine (i.e., privilege mode .. etc). <u>About Implementing Instructions:</u>

- The implementation must use these accessor functions to read and write to registers, memory and set/reset, test the status bits.
- The implementation must use the logic operations defined in logic.c to implement all functionality.
  - Note that the function lc3\_evaluate\_instruction() follows more or less the description given in Lectures 10 and 11.
- You should refer to Appendix A of the textbook to understand how to implement the individual instructions.

#### For example:

Look at the description for the AND instruction (from page 527 in Appendix A) from the textbook.

Operation
if (bit[5] == 0)
 DR = SR1 AND SR2;
else
 DR = SR1 AND SEXT(imm5);
setcc();

Here it says that if bit-5 of the instruction is set, then SR1 and SR2 are selected, otherwise SR1 and the sign extended register are selected. The sign extender must operate on the last five bits of the instruction. In the provide implementation, a function called sign\_extend\_imm5() performs this operation. Now look at the implementation for AND:

```
lc3_word_t ir_ = (lc3_word_t)read_register(IR);
lc3_register_names_t dr_ = gpr_select(ir_.bit[4], ir_.bit[5], ir_.bit[6]);
lc3_register_names_t sr1_ = gpr_select(ir_.bit[7], ir_.bit[8], ir_.bit[9]);
```

First, the destination register and source register-1 are selected, and we get the contents of IR.

if (ir\_.bit[10] == 0) {

Next, the instruction register's bit 5 is tested and the appropriate operation is selected. To do an immediate add:

The function to sign extend the immediate 5-bit value is called. This function implicitly writes to the SEXT register. Then the two words are ANDed using word\_and, and then finally the resulting word is stored back in the destination register. Note that word\_add() accepts its parameters as lc3\_word\_t struct variables. Therefore we needed to use temporary variables.

## **Submission Instructions**

Submit a tar.gz file called PA6.tar.gz with your implementation.

If you were unable to complete Part-A (i.e., tokenizing and parsing the symbol table file) and end up using infilesparser.o.PROVIDED, then also submit a README file describing your progress with implementing Part-A to be considered for partial credit. Do the same if you were unable to complete Part-B.

#### Failure to submit the README file will result in <u>no partial credit</u>.

#### **Grading Criteria**

This assignment will be graded in full for having the required functionality. Non-functional programs will only receive partial credit. We will award 30% for implementing Part-A, and the remaining 70% of the points will be awarded for a fully working Part-B implementation.

Note that in order to receive partial credit for non-functional programs, your code must be well documented and well structured, otherwise the partial credit will be minimal (and remember, you will not receive any partial credit if you fail to submit the README file).

#### Late Policy

Late submissions are not allowed for this assignment.