

Mathematical Induction

Rosen Chapter 5



Motivation

➤ Show that any postage of $\geq 8\text{¢}$ can be obtained using 3¢ and 5¢ stamps.

➤ First check for a few values:

$$8\text{¢} = 3\text{¢} + 5\text{¢}$$

$$9\text{¢} = 3\text{¢} + 3\text{¢} + 3\text{¢}$$

$$10\text{¢} = 5\text{¢} + 5\text{¢}$$

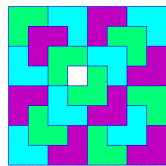
$$11\text{¢} = 5\text{¢} + 3\text{¢} + 3\text{¢}$$

$$12\text{¢} = 3\text{¢} + 3\text{¢} + 3\text{¢} + 3\text{¢}$$

➤ How to generalize this?

Motivation

- Let n be a positive integer. Show that every $2^n \times 2^n$ chessboard with one square removed can be tiled using triominoes, each covering three squares at a time.




Motivation

- Prove that for every positive value of n ,
 $1 + 2 + \dots + n = n(n + 1)/2$.

Motivation

- Many mathematical statements have the form:
 $\forall n \in \mathbb{N}, P(n)$ $P(n)$: Logical predicate
- Example: For every positive value of n ,
 $1 + 2 + \dots + n = n(n + 1)/2$.
- Predicate – propositional function that depends on a variable, and has a truth value once the variable is assigned a value.
- Mathematical induction is a proof technique for proving such statements






Proving P(3)

- Suppose we know: $P(1)$ and $P(n) \rightarrow P(n + 1) \forall n \geq 1$.
 Prove: $P(3)$
 - Proof:
 - $P(1)$. [premise]
 - $P(1) \rightarrow P(2)$. [specialization of premise]
 - $P(2)$. [step 1, 2, & modus ponens]
 - $P(2) \rightarrow P(3)$. [specialization of premise]
 - $P(3)$. [step 3, 4, & modus ponens]
- 
- We can construct a proof for every finite value of n**
- Modus ponens: if p and $p \rightarrow q$ then q

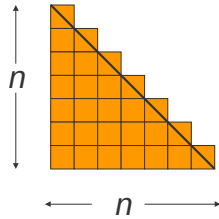
Example: $1 + 2 + \dots + n = n(n + 1)/2$.

- $P(n)$: $1+2+\dots+n = n(n + 1)/2$
- Prove: $P(n=1)$: $1 =? 1(1 + 1)/2 = 1$.
- Assume: $P(n=k)$ is true
 $1+2+\dots+k = k(k + 1)/2$
 - This is called the *Inductive Hypothesis*
- Use that assumption to prove $P(n=k+1)$ is true.
 $1+2+\dots+k+k+1 =? (k + 1)(k + 2)/2$.
 $(1+2+\dots+k) + k+1 =? (k+1)(k+2)/2$
 $k(k+1)/2 + k+1 =? (k+1)(k+2)/2$ [by Ind. Hyp.]
 $k(k+1)/2 + 2(k+1)/2 =? (k+1)(k+2)/2$
 $(k^2+k+2k+2)/2 =? (k+1)(k+2)/2$
 $(k^2+k+2k+2)/2 = (k^2 + k + 2k + 2)/2$

A Geometrical interpretation

- 1: 
- 2: 
- 3: 
- Put these blocks, which represent numbers, together to form sums:
- $1 + 2 =$ 
- $1 + 2 + 3 =$ 

A Geometrical interpretation



Area is $n^2/2 + n/2 = n(n + 1)/2$

The Principle of Mathematical Induction

- Let $P(n)$ be a statement that, for each natural number n , is either true or false.
- To prove that $\forall n \in \mathbb{N}, P(n)$, it suffices to prove:
 - $P(n=1)$ is true. (or sometimes some other n)
(base case)
 - $\forall n \in \mathbb{N}, P(n) \rightarrow P(n + 1)$. (inductive **step**)
- This is not magic.
- It is a recipe for constructing a proof for an arbitrary $n \in \mathbb{N}$.

Mathematical Induction and the Domino Principle

If

the 1st domino falls over

and

the n th domino falls over **implies** that domino $(n + 1)$ falls over

then

domino n falls over for all $n \in \mathbb{N}$.

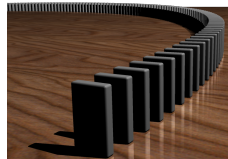


image from <http://en.wikipedia.org/wiki/File:Dominoeffect.png>

Proof by induction

- 3 steps:
 - Prove $P(n=1)$. [the **basis**]
 - Assume $P(n=k)$ [the **induction hypothesis**]
 - Using $P(n=k)$, prove $P(n=k+1)$ [the **inductive step**]

Example

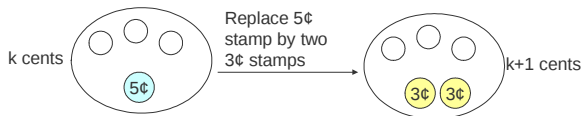
- Show that any postage of $\geq 8\text{¢}$ can be obtained using 3¢ and 5¢ stamps.
- First check for a few values:
 - $8\text{¢} = 3\text{¢} + 5\text{¢}$
 - $9\text{¢} = 3\text{¢} + 3\text{¢} + 3\text{¢}$
 - $10\text{¢} = 5\text{¢} + 5\text{¢}$
 - $11\text{¢} = 5\text{¢} + 3\text{¢} + 3\text{¢}$
 - $12\text{¢} = 3\text{¢} + 3\text{¢} + 3\text{¢} + 3\text{¢}$
- How to generalize this?

Example

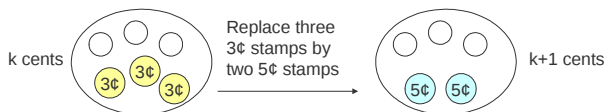
- Let $P(n)$ be the statement "n cents postage can be obtained using 3¢ and 5¢ stamps".
- Want to show that "P(n=k) is true" implies "P(n=k+1) is true" for all $k \geq 8$.
- How can you increase k by 1?
 - 2 cases:
 - 1) P(k) is true and the k cents contain at least one 5¢.
 - 2) P(k) is true and the k cents do not contain any 5¢.

Example

Case 1: k cents contain at least one 5¢ stamp.



Case 2: k cents do not contain any 5¢ stamp. Then there are at least three 3¢ stamp.



Examples

- Show that $1 + 2 + 2^2 + \dots + 2^n = 2^{n+1} - 1$
- Show that for $n \geq 4$ $2^n < n!$
- Show that $n^3 - n$ is divisible by 3 for every positive n.
- Show that $1 + 3 + 5 + \dots + (2n+1) = (n+1)^2$
- Prove that a set with n elements has 2^n subsets

All horses have the same color



- *Base case:* If there is only one horse, there is only one color.
- *Induction step:* Assume as induction hypothesis that within any set of n horses, there is only one color. Now look at any set of $n + 1$ horses. Number them: $1, 2, 3, \dots, n, n + 1$. Consider the sets $\{1, 2, 3, \dots, n\}$ and $\{2, 3, 4, \dots, n + 1\}$. Each is a set of only n horses, therefore within each there is only one color. But the two sets overlap, so there must be only one color among all $n + 1$ horses.
- This is clearly wrong, but can you find the flaw?

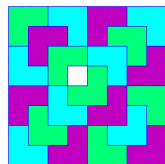
All horses have the same color



- The inductive step requires that $k \geq 2$, otherwise there is no intersection! So $P(2)$ should be the base case, which is obviously incorrect.
- In the book there is a similar example.

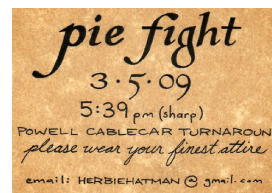
More induction examples

- Let n be a positive integer. Show that every $2^n \times 2^n$ chessboard with one square removed can be tiled using triominoes, each covering three squares at a time.



Odd Pie Fights

- An odd number of people stand at mutually distinct distances. Each person throws a pie at their nearest neighbor. Show that there is at least one survivor.



<http://laughingsquid.com/san-francisco-pie-fight-at-the-powell-street-cable-car-turnaround/>

Strong induction

- Induction:
 - $P(n=1)$ is true.
 - $\forall k \in \mathbb{N}, P(n=k) \rightarrow P(n=k+1)$.
 - Implies $\forall k \in \mathbb{N}, P(n=k)$
- Strong induction:
 - $P(n=1)$ is true.
 - $\forall k \in \mathbb{N}, (P(n=1) \wedge P(n=2) \wedge \dots \wedge P(n=k)) \rightarrow P(n=k+1)$
 - Implies $\forall k \in \mathbb{N}, P(n=k)$ is true.

Example

- Prove that all natural numbers ≥ 2 can be represented as a product of primes.
- **Basis:** 2: 2 is a prime.
- **Assume** that $1, 2, \dots, n$ can be represented as a product of primes.

Example

- **Show** that $n+1$ can be represented as a product of primes.
 - If $n+1$ is a prime: It can be represented as a product of 1 prime, itself.
 - If $n+1$ is composite: Then, $n + 1 = ab$, for some $a, b < n + 1$.
Therefore, $a = p_1 p_2 \dots p_k$ & $b = q_1 q_2 \dots q_l$, where all p_i and q_i are primes.
Represent $n+1 = p_1 p_2 \dots p_k q_1 q_2 \dots q_l$ which is a product of primes!

Breaking chocolate

Theorem: Breaking up a chocolate bar with n “squares” takes $n-1$ breaks.



Induction and Recursion

- Induction is useful for proving correctness of recursive algorithms

- Example

```
// Returns base ^ exponent.
// Precondition: exponent >= 0
public static int pow(int x, int n) {
    if (n == 0) {
        // base case; any number to 0th power is 1
        return 1;
    } else {
        // recursive case: x^n = x * x^(n-1)
        return x * pow(x, n-1);
    }
}
```

Induction and Recursion

```
public static int pow(int x, int n) {
    if (n == 0){
        return 1;
    } else {
        return x * pow(x, n-1);
    }
}
```

Claim: the algorithm correctly computes x^n .

Proof: By induction on n

Base case: $n = 0$: it correctly returns 1

Inductive step: assume that for n the algorithm correctly returns x^n .

Then for $n+1$ it returns $x * x^n = x^{n+1}$.

Induction and Recursion

- $n!$ of some integer n can be characterized as:
 $n! = 1$ for $n = 0$; otherwise
 $n! = n(n-1)(n-2) \dots 1$
- Want to write a recursive method for computing it. We notice that $n! = n(n-1)!$
- This is all we need to put together the method:

```
public static int factorial(int n) {
    if (n == 0) {
        return 1;
    } else {
        return n * factorial(n-1);
    }
}
```

Induction in CS

- Induction is a powerful tool for showing algorithm correctness – not just for recursive algorithms (CS320)