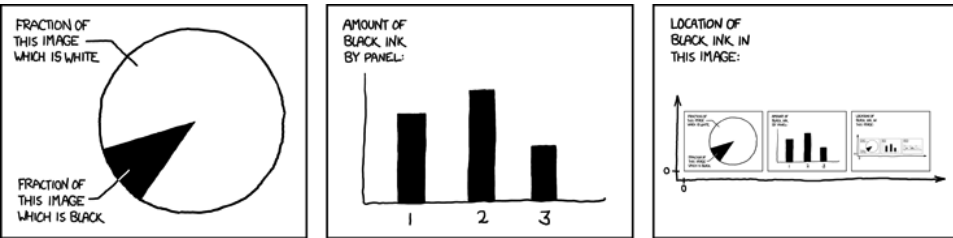


## More Recursion!



<http://xkcd.com/688/>

<http://xkcd.com/981/>

## Recursion - examples

- Problem: given a string as input, return the string with characters reversed.
- Base case?
- Recursion

## Tail recursion

- Tail recursion is a recursive call that occurs as the last action in a method.
- This is not tail recursion:

```
public int factorial(int n){
    if (n==0)
        return 1;
    return n * factorial(n-1);
}
```

- How can we make the call to factorial the last thing?

## Tail recursion

- Tail recursion is a recursive call that occurs as the last action in a method.
- This is not tail recursion:

```
public int factorial(int n){
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    return n * factorial(n-1);
}
```

- How can we make the call to factorial the last thing?
- Yep! Must use \* in a new argument.

## Tail recursion

- Non tail-recursive:

```
public int factorial(int n){
    if (n == 0)
        return 1;
    return n * factorial(n-1);
}
```

- Tail-recursive:

```
public int factorial(int n, int product) {
    if (n == 0)
        return product;
    return factorial(n-1, n * product);
}
```

## Dictionary lookup

- Suppose you're looking up a word in the dictionary (paper one, not online!)
- You probably won't scan linearly through the pages – inefficient.
- What would be your strategy?

## Tail recursion

- Let's hide this additional argument:

```
public int factorial(int n) {
    return factorialTail(n, 1);
}
private int factorialTail(int n, int product) {
    if(n == 0)
        return product;
    return factorialTail(n-1, n * product);
}
```

- But why would you care? Compilers can optimize memory usage when they detect tail recursion. When making a recursive call, you no longer need to save the information about the local variables within the calling method.

## Binary search

```
binarySearch(dictionary, word){

    // base case
    ???

    else { // recursive case

        open the dictionary to a point near the middle
        determine which half of the dictionary contains word

        if (word is in first half of the dictionary) {
            binarySearch(first half of dictionary, word)
        }
        else {
            binarySearch(second half of dictionary, word)
        }
    }
}
```

# Binary search

```
binarySearch(dictionary, word) {  
  
    if (dictionary has one page) { // base case  
        scan the page for word  
    }  
  
    else { // recursive case  
  
        open the dictionary to a point near the middle  
        determine which half of the dictionary contains word  
  
        if (word is in first half of the dictionary) {  
            binarySearch(first half of dictionary, word)  
        }  
        else {  
            binarySearch(second half of dictionary, word)  
        }  
    }  
}
```

# Binary search

- Let's write a method called `binarySearch` that accepts a **sorted** array of integers and a target integer and returns the index of an occurrence of that value in the array.
  - If the target value is not found, return -1

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
value	-	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92

```
int index = binarySearch(data, 42); // 10  
int index2 = binarySearch(data, 66); // -1
```

# Binary search

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92

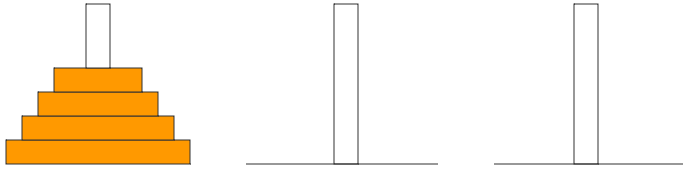
- How can we implement this?
  - Create two smaller arrays?
  - Pass start and end indicies?

# Binary search

```
// Precondition: a is sorted  
// Postcondition: Returns the index of an occurrence of the given value, or -1.  
public int binarySearch(int[] a, int target) {  
    return binarySearch(a, target, 0, a.length - 1);  
}  
  
// Recursive helper to implement search.  
private int binarySearch(int[] a, int target, int first, int last) {  
    if (first > last) {  
        return -1; // not found  
    } else {  
        int mid = (first + last) / 2;  
        if (a[mid] == target) {  
            return mid;  
        }  
  
        } else if (a[mid] < target) {  
            return binarySearch(a, target, mid+1, last);  
        } else {  
            return binarySearch(a, target, first, mid-1);  
        }  
    }  
}
```

# Towers of Hanoi

Example: Towers of Hanoi, move all disks to third peg without ever placing a larger disk on a smaller one.



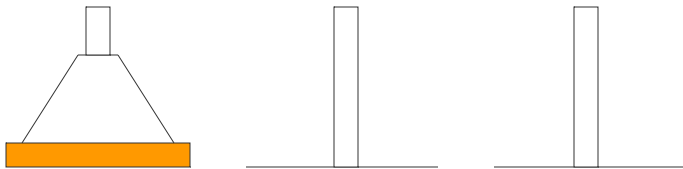
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# Try to find the pattern by cases

- One disk is easy
- Two disks...
- Three disks...
- Four disk...

# Towers of Hanoi

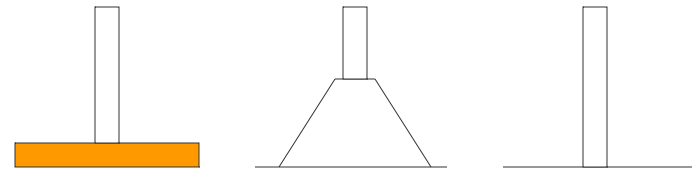
Example: Towers of Hanoi, move all disks to third peg without ever placing a larger disk on a smaller one.



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# Towers of Hanoi

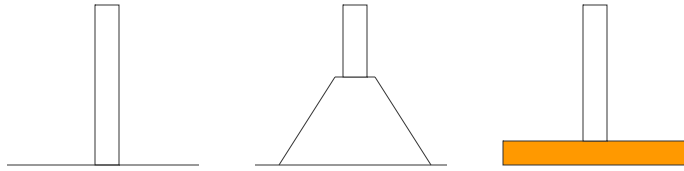
Example: Towers of Hanoi, move all disks to third peg without ever placing a larger disk on a smaller one.



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# Towers of Hanoi

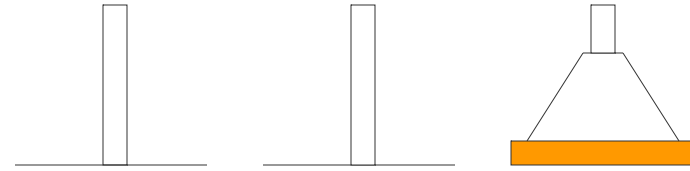
Example: Towers of Hanoi, move all disks to third peg without ever placing a larger disk on a smaller one.



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# Towers of Hanoi

Example: Towers of Hanoi, move all disks to third peg without ever placing a larger disk on a smaller one.



Let's go play with it at:

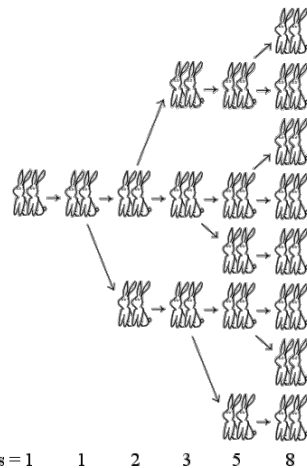
<https://www.mathsisfun.com/games/towerofhanoi.html>

[https://www.youtube.com/watch?v=4\\_KtPENqCb0](https://www.youtube.com/watch?v=4_KtPENqCb0)

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# Fibonacci's Rabbits

- Suppose a newly-born pair of rabbits, one male, one female, are put on an island.
  - A pair of rabbits doesn't breed until 2 months old.
  - Thereafter each pair produces another pair each month
  - Rabbits never die.
- How many pairs will there be after  $n$  months?



pairs = 1    1    2    3    5    8

image from: <http://www.jimloy.com/algebra/fibo.htm>

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# Fibonacci numbers

- The *Fibonacci numbers* are a sequence of numbers  $F_0, F_1, \dots, F_n$  defined by:

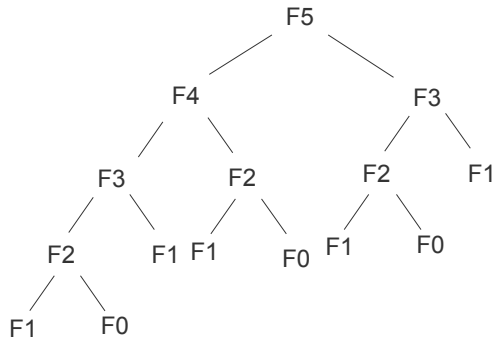
$$F_0 = F_1 = 1$$

$$F_i = F_{i-1} + F_{i-2} \text{ for any } i > 1$$

- Write a method that, when given an integer  $i$ , computes the  $n$ th Fibonacci number.

## Fibonacci numbers

- Let's run it for  $n = 1, 2, 3, \dots, 10, \dots, 20, \dots$
- If  $n$  is large the computation takes a long time! Why?



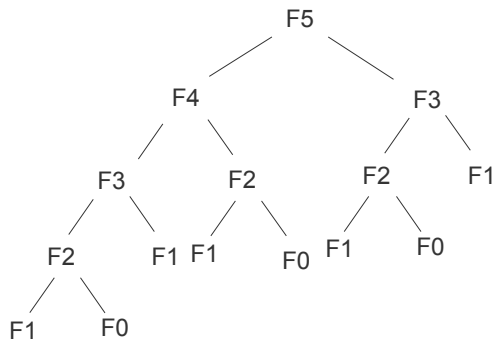
## Fibonacci numbers

- recursive Fibonacci was expensive because it made many, recursive calls
  - fibonacci( $n$ ) recomputed fibonacci( $n-1$ ), ..., fibonacci(1) many times in finding its answer!
  - This is a case where the sub-tasks handled by the recursion are redundant with each other and get recomputed

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## Fibonacci numbers

- Every time  $n$  is incremented by 2, the call tree more than doubles.



## Growth of rabbit population

1 1 2 3 5 8 13 21 34 ...

The fibonacci numbers themselves also grow rapidly: every 2 months the population at least **DOUBLES**

# Fractals - the Koch curve and Sierpinski Triangle

